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Solutions to Homework 6

- Number of operations for standard LU:  $\approx \frac{2}{3}n^3$  (SH, p. 57)  
 for Cholesky:  $\approx \frac{1}{3}n^3$  (SH, p. 95)

For each step of CG we have to compute

$\tau_k$ :  $n^2 + O(n)$  operations

$\beta_{k+1}$ :  $n^2$  operations, but can be reused from previous step

$\rho_k$ :  $O(n)$  operations once  $\beta_{k+1}$  is known

$\beta_k$ :  $n^2$  operations

$\alpha_k$ :  $O(n)$  operations once  $\beta_k$  is known

$x_{k+1}$ :  $O(n)$  operations

Thus, an efficient implementation will have  $2n^2 + O(n)$  operations per step, so  $n$  steps take  $2n^3 + O(n^2)$  ops; 6 times longer than Cholesky.

Note: In all this, it is common practice to count 1 addition and one multiplication as a single operation.

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3. a)  $L_0(x) = \frac{(x-x_1)(x-x_2)}{(x_0-x_1)(x_0-x_2)}$

$L_1(x) = \frac{(x-x_0)(x-x_2)}{(x_1-x_0)(x_1-x_2)}$

$L_2(x) = \frac{(x-x_0)(x-x_1)}{(x_2-x_0)(x_2-x_1)}$

$p(x) = L_0(x) \underbrace{f(x_0)}_{=: y_0} + L_1(x) \underbrace{f(x_1)}_{=: y_1} + L_2(x) \underbrace{f(x_2)}_{=: y_2}$

(Further simplifications do not gain much.)

b)  $f'(x) \approx p'(x) = \frac{(x-x_1)(x-x_2)}{(x_0-x_1)(x_0-x_2)} y_0 + \frac{(x-x_0)(x-x_2)}{(x_1-x_0)(x_1-x_2)} y_1 + \frac{(x-x_0)(x-x_1)}{(x_2-x_0)(x_2-x_1)} y_2$

$\Rightarrow f'(x) \approx \frac{(x_1-x_2)}{(x_0-x_1)(x_0-x_2)} y_0 + \left( \frac{1}{x_1-x_2} + \frac{1}{x_1-x_0} \right) y_1 + \frac{x_1-x_0}{(x_2-x_0)(x_2-x_1)} y_2$

$f''(x) \approx p''(x) = 2 \left[ \frac{y_0}{(x_0-x_1)(x_0-x_2)} + \frac{y_1}{(x_1-x_0)(x_1-x_2)} + \frac{y_2}{(x_2-x_0)(x_2-x_1)} \right]$

c) Equidistant nodes:  $x_1 - x_0 = h = x_2 - x_1$ ;  $x_2 - x_0 = 2h$

$\Rightarrow f'(x_1) \approx \frac{-h}{(-h)(-2h)} y_0 + \left( \frac{1}{-h} + \frac{1}{h} \right) y_1 + \frac{h}{2h \cdot h} y_2$

$= \frac{y_2 - y_0}{2h}$

$f''(x_1) \approx 2 \left[ \frac{y_0}{(-h)(-2h)} + \frac{y_1}{h \cdot (-h)} + \frac{y_2}{2h \cdot h} \right] = \frac{y_0 - 2y_1 + y_2}{h^2}$

(3)

4. It is easy to see that  $p(x)$  is a polynomial of degree at most  $n+1$ . Therefore, if  $p(x)$  interpolates all  $n+2$  points, it is the Lagrange polynomial (the unique polynomial of degree at most  $n+1$  with this property).

$$\text{(check: } \bullet p(x_0) = \frac{(x_0 - x_0) + (x_0 - x_{n+1}) q(x_0)}{x_{n+1} - x_0}$$

$$= q(x_0) = y_0 \quad \checkmark \quad (\text{as } q \text{ interpolates at } x_0)$$

$$\bullet p(x_{n+1}) = \frac{(x_{n+1} - x_0) + (x_{n+1} - (x_{n+1} - x_{n+1})) q(x_{n+1})}{x_{n+1} - x_0}$$

$$= + (x_{n+1}) = y_{n+1} \quad \checkmark \quad (\text{as } + \text{ interpolates at } x_{n+1})$$

• For  $j = 1, \dots, n$ :  $+ (x_j) = q(x_j) = y_j$ , so that

$$p(x_j) = \frac{(x_j - x_0) + (x_j - (x_j - x_{n+1})) q(x_j)}{x_{n+1} - x_0}$$

$$= \frac{(x_j - x_0) y_j - (x_j - x_{n+1}) y_j}{x_{n+1} - x_0}$$

$$= y_j \quad \checkmark$$