

1. Show that in a finite graph without cycles, there is a vertex of valency at most one. (10)

Solution 1: Let $\langle v_1, \dots, v_k \rangle$ be a chain of maximal length.

(Which exists since the graph is finite.) Then v_k is a vertex of valency at most one, for if not, the chain would either not be maximal, or there would have to be a cycle.

Solution 2: Take Euler's formula applied to a connected component of the graph. Since there are no cycles, $|F| = 1$, so

$$|V| - |E| = 1$$

$$\text{But } \sum_{v \in G} s(v) = 2|E| = 2|V| - 2$$

So there must be at least two vertices with valency no more than 1.

2. Let l be a fixed line in the plane. Recall that a *glide reflection* with axis l is a transformation $U = R_l \Pi$ where R_l is the line reflection about l and Π is some nonidentity translation which leaves l invariant.

- Show that $R_l \Pi = \Pi R_l$.
- Show that $U^{-1} = R_l \Pi^{-1}$.
- Show that U^{-1} is a glide reflection with axis l .
- Consider the set of all glide reflections with axis l . Is it a group? If not, describe the group it generates.

(5+5+5+5)

(a) WLOG let l be the x -axis and let Π be right translation by one unit. Then

$$\begin{pmatrix} x \\ y \end{pmatrix} \xrightarrow{\Pi} \begin{pmatrix} x+1 \\ y \end{pmatrix} \quad \text{and} \quad \begin{pmatrix} x \\ y \end{pmatrix} \xrightarrow{R_l} \begin{pmatrix} x \\ -y \end{pmatrix}$$

It is then trivial to check that

$$\begin{pmatrix} x \\ y \end{pmatrix} \xrightarrow{R_l \Pi} \begin{pmatrix} x+1 \\ -y \end{pmatrix} \quad \text{and so does } \Pi R_l$$

(A geometric argument, looking at the image of a point under the various maps is equally easy.)

$$(b) \quad R_l \Pi^{-1} U = R_l \Pi^{-1} R_l \Pi \stackrel{(a)}{=} R_l \Pi^{-1} \Pi R_l = R_l^2 = \text{Id}$$

$$\Rightarrow R_l \Pi^{-1} = U^{-1}.$$

(c) Since Π^{-1} is also a non-identity translation which leaves l invariant, by (b), $U^{-1} = R_l \Pi^{-1}$ satisfies the definition of a glide reflection.

(d) The set of glide reflections with axis l does not contain the identity, so is not a group.

The group generated by it contains all translations along l (the composition of two glide reflections), hence also reflection about l ($R_l = U \circ \Pi^{-1}$!).

Thus, it is homeomorphic to the group $\mathbb{R} \times \mathbb{Z}_2$.

(6)

$$S_2 M = \begin{pmatrix} 0 & \dots & 0^{m_N} \\ m_1 & & 0 \\ \vdots & \ddots & \vdots \\ 0 & \dots & 0^{m_{N-1}} \end{pmatrix}$$

$$(S_2 M)^2 = \begin{pmatrix} 0 & \dots & 0^{m_{N-1} m_N} & 0 \\ m_1 m_2 & & & m_N m_1 \\ \vdots & \ddots & & \vdots \\ 0 & \dots & 0 & 0 \\ m_{N-2} m_{N-1} & & 0 & 0 \end{pmatrix}$$

$$\vdots$$

$$(S_2 M)^N = \begin{pmatrix} m_1 m_2 \dots m_N & & 0 \\ & \ddots & \\ 0 & & m_1 m_2 \dots m_N \end{pmatrix} = \begin{cases} I & \text{if } m_1 \dots m_N = 1 \\ -I & \text{if } m_1 \dots m_N = -1 \end{cases}$$

So if $m_1 \dots m_N = 1$, the set $G = \{(S_2 M)^i : i \in \mathbb{Z}\}$ contains N elements, if $m_1 \dots m_N = -1$, it contains $2N$ elements. Further, $I \in G$.

By the rules of matrix exponentiation, identifying the element $(S_2 M)^i \in G$ with $i \in \mathbb{Z}_N$ (resp. \mathbb{Z}_{2N}), we see that this defines a group homomorphism with \mathbb{Z}_N (resp. \mathbb{Z}_{2N}), hence G is a group.

(Note: it is straightforward, but tedious, to verify the group axioms one-by-one.)

(c) Let $X_i(t)$ denote the color of the i -th site, with 1 for black and -1 for white, and let $m_i = -1$ iff the i -th edge carries a marker. Then the evolution of the ring is given by

$$\begin{pmatrix} X_1(t) \\ \vdots \\ X_N(t) \end{pmatrix} = (S_2 M)^t \begin{pmatrix} X_1(0) \\ \vdots \\ X_N(0) \end{pmatrix},$$

i.e., G could be called the "Kac ring group".

(d) If $m_1 = \dots = m_N = -1$, then $S_2 M$ represents a (cyclic) shift and reflection. Moreover, $S_2 M$ satisfies the same commutator relation as a glide reflection, so it could be seen as a discrete glide reflection.

In this sense, the Kac ring evolution is a generalized glide reflection where some sites are reflected and others are not.

4. Solve the linear programming problem

$$\text{minimize } z = x + 3y$$

subject to

$$x + 2y \geq 2, \quad (A)$$

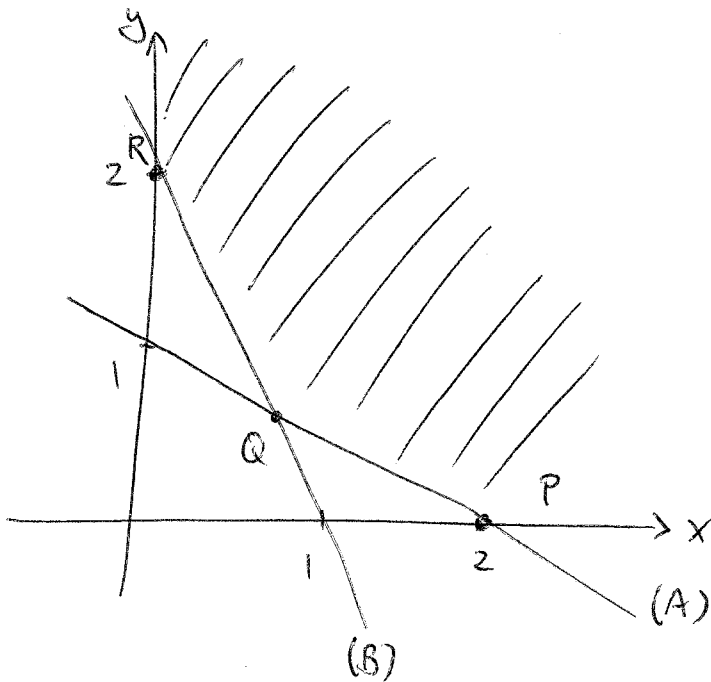
$$2x + y \geq 2, \quad (B)$$

$$x \geq 0,$$

$$y \geq 0,$$

using either the graphical method or the simplex method.

(10)



Compute coordinates of Q:

$$\begin{pmatrix} 1 & 2 & | & 2 \\ 2 & 1 & | & 2 \end{pmatrix} \rightarrow \begin{pmatrix} 1 & 2 & | & 2 \\ 0 & -3 & | & -2 \end{pmatrix}$$

$$\rightarrow \begin{pmatrix} 1 & 0 & | & \frac{2}{3} \\ 0 & 1 & | & \frac{2}{3} \end{pmatrix}$$

$$\Rightarrow Q = \left(\frac{2}{3}, \frac{2}{3}\right), \quad z = \frac{2}{3} + 2$$

$$P = (2, 0), \quad z = 2$$

$$R = (0, 2), \quad z = 6$$

\Rightarrow The minimum is attained at P with $z=2$.

5. The linear programming problem

$$\begin{aligned} & \text{maximize } \mathbf{c}^T \mathbf{x} \\ & \text{subject to } A\mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq 0 \end{aligned} \tag{P}$$

has a corresponding symmetric dual problem

$$\begin{aligned} & \text{minimize } \mathbf{b}^T \mathbf{y} \\ & \text{subject to } A^T \mathbf{y} \geq \mathbf{c}, \mathbf{y} \geq 0. \end{aligned} \tag{D}$$

Suppose that \mathbf{x} is feasible for (P) and \mathbf{y} is feasible for (D).

(a) Show that $\mathbf{c}^T \mathbf{x} \leq \mathbf{b}^T \mathbf{y}$.

(b) Conclude that if $\mathbf{c}^T \mathbf{x} = \mathbf{b}^T \mathbf{y}$, then \mathbf{x} and \mathbf{y} are optimal for their respective linear programming problems.

(5+5)

$$(a) \quad \mathbf{c}^T \mathbf{x} \stackrel{(D)}{\leq} (A^T \mathbf{y})^T \mathbf{x} = \mathbf{y}^T A \mathbf{x} \stackrel{(P)}{\leq} \mathbf{y}^T \mathbf{b} = \mathbf{b}^T \mathbf{y}$$

(b) Since $\mathbf{b}^T \mathbf{y}$ is an upper bound for the objective function of (P), attainment of the upper bound implies optimality for (P). Likewise for (D).

6. Recall that for $v \in \mathbb{C}^N$, the discrete Fourier transform of v is defined

$$\hat{v}_k = \frac{1}{N} \sum_{j=0}^{N-1} e^{-ijkh} v_j$$

with $h = 2\pi/N$ and for $k = 0, \dots, N-1$.

- (a) Set $w_j = e^{ijmh} v_j$ for $j = 0, \dots, N-1$. Express the discrete Fourier transform of w in terms of the discrete Fourier transform of v .
- (b) Let $v \in \mathbb{R}^N$ be a vector of *real* numbers. Show that its discrete Fourier transform satisfies

$$\overline{\hat{v}_k} = \hat{v}_{N-k},$$

where the overbar denotes the complex conjugate.

(5+5)

$$(a) \quad \hat{w}_k = \frac{1}{N} \sum_{j=0}^{N-1} \underbrace{e^{-ijkh} e^{ijmh}}_{= e^{-ij(k-m)h}} v_j = \hat{v}_{k-m}$$

$$(b) \quad \overline{\hat{v}_k} = \overline{\frac{1}{N} \sum_{j=0}^{N-1} e^{-ijkh} v_j} = \frac{1}{N} \sum_{j=0}^{N-1} e^{ijkh} v_j = \hat{v}_{-k}$$

Since \hat{v}_k is periodic with period N , we can write

$$\hat{v}_{-k} = \hat{v}_{N-k}$$

to translate back to the standard range of wave numbers.